

### **Task Farming in Contrail**

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### The Contrail Project



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# ConPaaS

- Contrail's Platform as a Service
  - PHP-based Web applications
  - MySQL
  - MapReduce
  - Task Farming
  - XtreemFS files system
  - Accessible via a common Web GUI





### **ConPaaS GUI**



### **ConPaaS Service Architecture** Cloud users Standard Web VM images interface Today: Service controller Task farming VM instance service Worker VM instances ConPaaS service contrail-project.eu

# **Task Farming**

- Dominant application type in grids
  - over 75% of all submitted tasks
  - over 90% of the total CPU-time consumption
  - [losup,Epema et al.]
- High-throughput applications (Condor style)
  - Parameter sweep
- Traditional execution model "grab and run"
  - Get as many machines as possible
  - Computation for free, best-effort execution
  - Desktop grids, clusters, ...



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### The promise of the cloud



 Elastic computing, get exactly the machines you need, exactly when you need them...

•Well, did we mention you have to pay for the hour?



## "Quality of Service"

### •Small Instance, \$0.085 per hour

- I.7 GB of memory, I EC2 Compute Unit (ECU)
- High-memory extra large, \$0.50 per hour
   17.1 GB memory, 6.5 ECU

### • High CPU medium, \$0.17 per hour

• I.7 GB of memory, 5 EC2 Compute Units

### Which one is faster for <u>my</u> application???

### Which one is cost efficient???

# **Bag Characteristics**

- Many independent tasks
  - All tasks are always ready to run
- Runtimes are unknown to the user
- Tasks have some (unknown) runtime distribution
- Simplifications:
  - Tasks can be aborted/restarted
  - No costs of input/output files (ongoing work)
  - No disruptive performance changes across clouds (e.g., with cache sizes that delay some tasks but not the others)



### **Cloud Characteristics**



- A cloud offering provides machines of certain properties like CPU speed and memory
  - All machines in a cloud offering are homogeneous
  - There is an upper limit of machines per cloud that a user can get
- A machine is charged per Accountable Time Unit (ATU);
   I hour, for example
- We call a cloud offering (machine type, price, max. number) a <u>cluster</u>
  - We are HPC guys, after all...



### What's the (scheduling) problem?

We are on a budget.We know nothing.



• We want to:

- Run all tasks from our bag on (cloud) clusters, without spending more than our budget
- Allocate/release machines dynamically while learning how fast our tasks execute on the different clusters
- If we learn that our budget is too low, give up
- Minimize makespan of the whole bag, if we can make it within budget



### BaTS: Budget-aware task scheduler

- Self scheduling tasks
- Reconfiguring cluster configurations



# The BaTS Story

- "Every good story has a beginning, a middle part, and an end."
- With BaTS:
  - Runtime and budget estimation
  - Throughput phase
  - Tail phase





### **Runtime Estimation**

- Statistics for sampling with replacement:
  - Bag of tasks can be described with pretty good accuracy from a small sample
  - We collect average and variance



## **Runtime Estimation**

- For each cluster (cloud machine type) we need a sample of +/- 30 completed tasks
  - (drawn at random)



# **Compact Sampling**

Assume: g(x) = a \* f(x)+b

Linear Regression: Replicate 7 tasks

Distribute rest of sample (30-7=23) over all clusters

Map samples to other clusters



### **Cluster Configuration**

 From the average speed of each cluster, (in tasks per minute) we can compute estimates for makespan (Te) and cost (Be) for a configuration from nodes of multiple clusters:

$$T_e = \frac{N}{\sum_{i=1}^{C_{nc}} \frac{a_i}{T_i}} \quad ; \quad B_e = \left\lceil \frac{T_e}{ATU} \right\rceil * \sum_{i=1}^{C_{nc}} a_i * c$$

We minimize Te while keeping Be <= B using a modified Bounded Knapsack Problem (BKP)
The BKP can be solved in pseudo-polynomial time, as a 0-1 knapsack problem via linear programming
BaTS chooses the configuration with minimal Te for Be <= B</li>

# **Budget Estimation**

- User must make the trade-off between cost and completion time
- BaTS provides the user with choice (cost, time), using cluster configurations computed from the sampling phase:
  - Cheapest makespan
  - Cheapest makespan +10/20% cost
  - Fastest makespan -10/20% cost
  - Fastest makespan
  - (more options are possible)
- Each configuration (in fact) consists of the numbers of machines per cluster



### **BaTS: Throughput Phase**

- Self scheduling tasks
- Reconfiguring cluster configurations regularly





# **Progress Monitoring**

- BaTS starts from the user-selected, initial configuration
- At regular intervals (e.g., 5 minutes), BaTS re-evaluates the configuration
  - 1. Update average and variance per cluster
  - 2. Re-evaluate the machine configuration
- Execution on real machines adds some complexity:
  - Individually requested from the cloud provider(s), startup time before being ready
  - Each machine has its own end of the next ATU



# Re-evaluate the machine configuration

- Solve the remaining problem
  - Less tasks
  - Less money left
  - Track already-paid time left on machines
- If budget violation expected, get more machines with better price/performance ratio, and drop others
- If makespan violation expected, get more fast machines, and drop others
- If both budget and makespan violations expected, call mummy the user



### Fluid vs.Discrete Models

- BaTS (the BKP solver) allocates machines per full ATU
- Assumes a "fluid" model of computing time



### Fluid vs.Discrete Models

- Tasks, however, are sequential, cannot be split across "leftover" cycles
- Tasks on machines in final ATU:



### The End is Near!

- The tail phase needs some special consideration
- Bags with high variance may overrun predicted makespan (and thus budget)
- Even without overrunning, towards the end machines remain idle





## **BaTS' Tail Phase**

- As soon as a machine can not be assigned a task, BaTS switches to the *tail phase:* 
  - Replicate running tasks onto idle machines
- Which task to replicate?
  - The one that will terminate last!
- OK, how do we know?
  - Estimate completion time based actual runtime:
    - "Task *i* is running for 12 minutes now, what is its expected completion time, given the observed average and variance of the bag?"
  - Estimate completion time onto the idle machine (starting from scratch)
  - If shorter, replicate



# BaTS' Tail Phase (2/2)

- Do we need to start earlier?
- In the throughput phase, the average runtime determines the speed.
  - According to the *central limit theorem*, this no longer holds, once the population is smaller than a threshold (the same as the sample size in the beginning, +/- 30)
- With the threshold reached, BaTS migrates tasks to faster machines.
  - Same as replication, but original task is killed.
  - This frees a slow machine for a hopefully shorter task.



## **BaTS' Tail Phase Evaluation**

- We compare the following options:
  - No tail phase optimization.
  - Stochastic replication (based on completion time prediction)
  - Replication with perfect knowledge (theoretical optimum)
  - Replication with random task selection (no knowledge)
  - Replication plus migration



# Types of Bags Used

- Normal distribution
- Truncated Levy distribution (heavy tailed)
- Multi-modal distribution (real world data)



## **Normal Distribution**

- Simulator runs
- 30 bags each
- 30 runs each





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### **Heavy-tailed Distribution**



## **Multi-modal Distribution**

50

-10 -20 -30 -40

makespan compared to estimation (in %)



# **Tail Phase Findings**

- Doing "nothing" is the only bad option
- Replication works fine
  - Even with random selection
    - But has higher error rate
- Additional migration seems not to be worth the effort
  - The price we pay (kill running task) seems to outweigh the benefits



### BaTS on the Amazon Spot Market

- So far, we used "on demand" instances
  - Fixed price per hour
- Amazon spot market:
  - Same (on demand) machine types at different prices
    - Users "bid" a price for a machine (of a type)
    - If the bid is >= the current spot price, user gets the machine
    - If the spot prices exceeds the bid, the user is kicked out without prior notice
      - (and is reimbursed for the aborted hour)



### Spot Market: pros and cons

- Pro:
  - We might get machines cheaper
  - In practice, spot prices hardly ever change (boring)
- Con:
  - Tasks might get aborted
    - (we also do this ourselves, no problem)
  - Total budget fluctuates
  - Getting a spot instance takes +/- 8 minutes (before the booting starts)



### **BaTS Sampling for the Spot Market**



### New research problem: What is a good bidding strategy for spot machines?



### **Bidding Strategies**

- Maximum price
  - Determine the max price at which a spot instance is more cost efficient than the most profitable on-demand instance:  $T_p$

$$Max_i = \frac{T_p}{T_i} * c_p - \varepsilon$$

- Current price
  - Always get spot instances, the cheapest option at the moment of execution
- Average price
  - Literally the average between "current" and "maximum", in between the two extremes



### **Spot Market Estimations**

- Using max. 10 instances each of t1.micro, m1.small, m1.medium
- Bag with 18000 tasks (average 32, 15, and 8 seconds)
- Max. bid used: \$0.02 for t1.micro, \$0.007 for m1.small
   and \$0.015 for m1.medium

No clear "winner" strategy. The user simply gets more options...









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# **Spot Market Findings**

- It is too early for final conclusions.
- Opens more choices for the cost-savvy user.
- The current implementation only uses the current spot prices (no history)
- Taking long-term spot prices into account, user might opt for a hard cost limit:
  - Place a low bid and wait until the price drops
  - Interrupt the whole computation if price goes up during the computation



## Conlusions

- BaTS gives the user control over and choice from several cloud offers
  - Run cheaper and longer
  - Or run faster with higher budget
- Learning stochastic properties of tasks works well in the absence of runtime estimates
- Next steps:
  - Fully integrate file I/O
  - Handle fluctuating node performance (ongoing)
  - Support workflows (tasks with dependencies)
  - Fault tolerance Resilience
  - Dig deeper into spot mark









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